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Outline



- Avoiding Squares in Partial Words
 Binary Alphabet
 - Ternary Alphabet
- 3 Avoiding Squares in Binary Cube-free Partial Words

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1 Preliminaries

2 Avoiding Squares in Partial Words

- Binary Alphabet
- Ternary Alphabet

8 Avoiding Squares in Binary Cube-free Partial Words

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Words and Partial Words

Alphabet: Σ – a non-empty finite set of symbols, or *letters* Hole: \diamond – a character outside the alphabet that is viewed as representing any letter of the alphabet.

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A word is a concatenation of symbols from Σ A partial word is a concatenation of symbols from $\Sigma_{\diamond} = \Sigma \cup \{\diamond\}.$ Σ^* denotes the set of words over Σ

Length of w: |w| – the number of symbols in w

Example

For example, in the binary alphabet $\{a, b\}$,

◇ : a, b
◇ : aa, ab, ba, bb
bb◇a◇ : bbaaa, bbaab, bbbaa, bbbab
abaa : abaa

Since each \diamond may represent any letter in the alphabet, a partial word with *h* holes *represents* $|\Sigma|^h$ words.

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Note that a (full) word is a concatenation of letters from Σ_\diamond with no holes.

Compatibility

For two partial words u, v, where |u| = |v|, u is *compatible* with v ($u \uparrow v$), if u and v agree at each position where neither has a hole.

In other words, u and v could potentially represent the same word.

For example,

ab⇔	a⇔b⇔	ab⇔⇒
a⇔b⇔	aab◇	aab⇔
ab⇔⇔ † a⇔b⇔	a◊b◊↑ aab◊	ab⇔≎†∕aab⇔

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ab⇔⇔ ↑ a⇔b⇔	a◊b◊↑ aab◊	ab⇔≎†∕aab⇔	

A word of the form $uu = u^2$ is a square (word) if it is nonempty. A word of the form $vvv = v^3$ is a cube (word) if it is nonempty.

A partial word W has a square S if S is compatible with some consecutive subword of W.

For example, over $\{a, b, n, s, t\}$,

bananas aaaaa banan\s

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A partial word W has a square S if S is compatible with some consecutive subword of W.

For example, over $\{a, b, n, s, t\}$, bananas aaaaa $banan \diamond s$ $anan = (an)^2$

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bananasaaaaabanan
$$\diamond$$
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Morphism

A morphism, $\varphi : A^* \to \Sigma^*$, is a function such that if $u = a_0 a_1 a_2 \cdots$ (where $a_i \in A$ for all i), $\varphi(u) = \varphi(a_0)\varphi(a_1)\varphi(a_2)\cdots$.

For example, let

arphi :	$\{a, b, c\}^*$	\rightarrow	$\{d, e\}^*$
	а	\mapsto	d
	Ь	\mapsto	ee
	С	\mapsto	de

Then $\varphi(abc) = \varphi(a)\varphi(b)\varphi(c) = deede$.

Note that the domain of the morphism is typically a (full) word.

A morphism, $\varphi : A^* \to \Sigma^*$, is *n*-uniform if $|\varphi(a)| = n$ for all $a \in A$. A morphism, $\varphi : A^* \to \Sigma^*$, is *iterative* if $A = \Sigma$.

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2 Avoiding Squares in Partial Words

- Binary Alphabet
- Ternary Alphabet

3 Avoiding Squares in Binary Cube-free Partial Words

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History

In 1906, 1912, Thue presents an infinite square-free word over a 4-letter alphabet, which he uses to construct an infinite square-free ternary word.

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ab

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aba

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In 1906, 1912, Thue presents an infinite square-free word over a 4-letter alphabet, which he uses to construct an infinite square-free ternary word.

In 1974, Entringer, Jackson, and Schatz give an infinite binary word that avoids all squares xx with $|x| \ge 3$, and prove the bound 3 is optimal.

In 1995, Fraenkel and Simpson give an infinite binary word that has only three squares a^2 , b^2 , $(ab)^2$.

In 2005, Rampersad, Shallit, and Wang simplify the proof for the result above using two uniform mophisms.

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In 2005, Rampersad, Shallit, and Wang simplify the proof for the result above using two uniform mophisms.

In 2007, Manea and Mercaş introduce the concept of squares in partial words.

In 2009, we investigate the avoidance of squares in both partial words and cube-free partial words.

Avoiding Squares in Partial Words

Binary Alphabet

Goal

Let g(h, s) denote the maximum length of a binary partial word with h holes that has at most s distinct squares.

The goal is to fill up the table below.



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Avoiding Squares in Partial Words

Binary Alphabet

Simple Results for Words (Franekel and Simpson)

g(0,1) = 7 (aaabaaa, abaaaba, abbbabb, their complements, reverses). g(0,2) = 18 (abaabbaaabbbaabbab, babbaabbbaaabbaaba). $g(0,s) = \infty$ for all integers $s \ge 3$.

Binary Alphabet

Simple Results for Partial Words

one hole :

 $g(1,0)=1 (\diamond), g(1,1)=5 (\diamond aaba), g(1,2)=16 (\diamond abbaaabbbaabbaab)$ two holes:

 $g(2,0){=}0, \qquad g(2,1){=}3 \ (\diamond a \diamond), \qquad g(2,2){=}14 \ (\diamond a b b a a a b b b a a b \diamond)$ three holes:

$$g(3,0)=0, \qquad g(3,1)=0, \qquad g(3,2)=9 \ (\diamond abb \diamond aab \diamond)$$

The only words valid for g(3,2) are $\diamond \diamond \diamond, \diamond a \diamond b \diamond, \diamond a \diamond b b a \diamond, \diamond a b b \diamond \diamond a a b \diamond$, their reverses and complements.

Moreover, for four holes there exist no partial word that has fewer than three squares.

Avoiding Squares in Partial Words

Binary Alphabet

Table so far

$s \setminus h$	0	1	2	3	4	5	6
0	3	1	0	0	0		
1	7	5	3	0	0		
2	18	16	14	9	0		
3	∞						
4							
5							

Avoiding Squares in Partial Words

Binary Alphabet

Rampersad, Shallit, and Wang's Morphisms, 1

5-letter 24-uniform morphism α :

а	\mapsto	abcdcbabcdeabcbabcdcbcde
b	\mapsto	abcbabcdedcdeabcdedcbcde
с	\mapsto	abcbabcdcbcdeabcdcbabcde
d	\mapsto	abcdcbcdedcdeabcdcbabcde
е	\mapsto	abcdcbcdeabcbabcdedcbcde

6-uniform morphism β :

а	\mapsto	abbbaa
b	\mapsto	babbaa
с	\mapsto	bbbaaa
d	\mapsto	bbaaba
е	\mapsto	bbaaab

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Avoiding Squares in Partial Words

Binary Alphabet

Rampersad, Shallit, and Wang's Morphisms, 1

5-letter 24-uniform morphism α :

- abcdcbabcdeabcbabcdcbcde \mapsto а h \mapsto abc babcdedcdeabcdedc bcde \mapsto abc babcdcbcdeabcdcbabcde С *→ abcdcbcdedcdeabcdcbabcde* d \mapsto abcdcbcdeabcbabcdedcbcde е

6-uniform morphism β :

а	\mapsto	abbbaa
b	\mapsto	babbaa
С	\mapsto	bbbaaa
d	\mapsto	bbaaba
е	\mapsto	bbaaab

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Binary Alphabet

No Holes

Lemma (Rampersad, Shallit, Wang)

If $u \in \{a, b, c, d, e\}^*$ is square-free and avoids the patterns ac, ad, ae, be, ca, da, eb, then $\alpha(u)$ is square-free and avoids the patterns ac, ad, ae, bd, be, ca, ce, da, db, eb, ec, aba, ede.

Lemma (Rampersad, Shallit, Wang)

If $u \in \{a, b, c, d, e\}^*$ is square-free and avoids the patterns ac, ad, ae, bd, be, ca, ce, da, db, eb, ec, edeaba, then the only squares in $\beta(u)$ are aa, bb, abab.

Theorem (Rampersad, Shallit, Wang)

There exists an infinite binary word, $w_0 = \beta(\alpha^{\omega}(a))$, that has only the three squares: a^2 , b^2 , $(ab)^2$.

Avoiding Squares in Partial Words

Binary Alphabet

One and Two Holes

Theorem (B.-S., C., M. (2011))

There exists an infinite binary partial word with one hole, $w_1 = \diamond \beta(\alpha^{\omega}(a))$, that has only the three squares: a^2 , b^2 , $(ab)^2$.

Theorem (B.-S., C., M. (2011))

There exists an infinite binary partial word with two holes, $w_2 = \Leftrightarrow \beta(\alpha^{\omega}(a))$, that has only the three squares: a^2 , b^2 , $(ab)^2$.

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Binary Alphabet

Proof of One Hole

Suppose not. Then, the 4th square must be compatible with $\diamond ua_0 u'$, where $\diamond ua_0 u'a_1$ is a prefix of $w_1 = \diamond \beta(\alpha^{\omega}(a))$ where $u = u' \in \{a, b\}^*$ and $a_0 \neq a_1$ are letters. Note that u = u' has prefix $\beta(ab) = abbbaababbaa$. Let |u| = 6q + r.

6-uniform morphism β :

- $egin{array}{ccc} a & \mapsto & abbbaa \ b & \mapsto & babbaa \end{array}$
- $c \mapsto bbbaaa$
- $d \mapsto bbaaba$
- $e \mapsto bbaaab$

If r=0, then a_0abbb is a prefix of $\beta(f)$ for some fIf r=1, then a_0abbb is a suffix of $\beta(f)$ for some fIf r=2, then a_0abb is a suffix of $\beta(f)$ for some fIf r=4, then bbbaab is $\beta(f)$ for some f

If r=5, since all images of β have unique 5-letter prefixes, $a_0 = a_1$.

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Binary Alphabet

Proof of One Hole

Suppose not. Then, the 4th square must be compatible with $\diamond ua_0 u'$, where $\diamond ua_0 u'a_1$ is a prefix of $w_1 = \diamond \beta(\alpha^{\omega}(a))$ where $u = u' \in \{a, b\}^*$ and $a_0 \neq a_1$ are letters. Note that u = u' has prefix $\beta(ab) = abbbaababbaa$. Let |u|=6q+r.

6-uniform morphism β :

~		abbbaa	$\beta(f)$ has suffix $a_0 ab \Rightarrow$	f = e
d	\mapsto	adddaa	$\beta(x)$ bhashs	~ ~ d
b	\mapsto	babbaa	$p(g) = bbaaba \Rightarrow$	g = a
с	\mapsto	bbbaaa	$eta(h)$ has prefix <i>bbaa</i> \Rightarrow	$h \in \{d, e\}$
d	\mapsto	bbaaba	If $h = d \alpha^{\omega}(a)$ has square dd	
е	\mapsto	bbaaab	If $h = e$, then $\alpha^{\omega}(a)$ has ede	

~ (**~**) .

If r = 3, then there are letters f, g, h where

Avoiding Squares in Partial Words

Binary Alphabet

Three and Four Holes

Proposition

For large m, $w_3 = \beta(\alpha^m(a)) \diamond$ has only the three squares: a^2 , b^2 , $(ab)^2$.

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For large m, $w_4 = \beta(\alpha^m(a)) \diamond bas$ only three squares: a^2 , b^2 , $(ab)^2$.

Theorem (B.-S., C., M. (2011))

There exist an arbitrarily long finite partial word with four holes, $\Leftrightarrow \beta(\alpha^m(a)) \Leftrightarrow$, that has only the three squares: a^2 , b^2 , $(ab)^2$.

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Avoiding Squares in Partial Words

Binary Alphabet

Table so far

$s \setminus h$	0	1	2	3	4	5	6
0	3	1	0	0	0		
1	7	5	3	0	0		
2	18	16	14	9	0		
3	∞	∞	∞	$<\infty?$	$<\infty?$		
4							
5							

Binary Alphabet

Results

We found that all binary partial words of the form $u_0 \diamond u_1$ with $|u_0| = |u_1| = 8$ have at least four distinct squares.

Thus, all partial words with at most three distinct squares and at least three holes must have the holes placed within the first, or last, eight positions.

Also, we found that all partial words with at least three holes within the first eight symbols that have no more than three distinct squares have maximum length 13.

Binary Alphabet

Results

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Theorem (B.-S., C., M., (2011))

All infinite partial words with at least three holes have at least four distinct squares.

Theorem (B.-S., C., M., (2011))

All partial words with more than four holes have at least four distinct squares.

Avoiding Squares in Partial Words

Binary Alphabet

Table so far

$s \setminus h$	0	1	2	3	4	5	6
0	3	1	0	0	0	0	
1	7	5	3	0	0	0	• • •
2	18	16	14	9	0	0	• • •
3	∞	∞	∞	$<\infty$	$<\infty$	0	• • •
4							
5							

Avoiding Squares in Partial Words

Binary Alphabet

Modification of the Morphism

5-letter 24-uniform morphism α :

а	\mapsto	abcdcbabcdeabcbabcdcbcde
b	\mapsto	abcbabcdedcdeabcdedcbcde
с	\mapsto	abcbabcdcbcdeabcdcbabcde
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 $e \mapsto abcdcbcdeabcbabcdedcbcde$

6-uniform morphism β :

- $\begin{array}{cccc} a & \mapsto & abbbaa \\ b & \mapsto & babbaa \\ c & \mapsto & bbbaaa \\ d & \mapsto & bbaaba \end{array}$
- $e \mapsto bbaaab$

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Avoiding Squares in Partial Words

Binary Alphabet

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b	\mapsto	abcbabcdedcdeabcdedcbcde
с	\mapsto	abcbabcdcbcdeabcdcbabcde
d	\mapsto	abcdcbcdedcdeabcdcbabcde

 $e \mapsto abcdcbcdeabcbabcdedcbcde$

6-uniform morphism β :

а	\mapsto	abbbaa
Ь	\mapsto	babbaa
с	\mapsto	bbbaaa
d	\mapsto	bbaaba
е	\mapsto	bbaaab

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Avoiding Squares in Partial Words

Binary Alphabet

Modification of the Morphism

5-letter 24-uniform morphism α :

а	\mapsto	abcdcbabcdeabcbabcdcbcde
b	\mapsto	abcbabcdedcdeabcdedcbcde
с	\mapsto	abcbabcdcbcdeabcdcbabcde
d	\mapsto	abcdcbcdedcdeabcdcbabcde

 $e \mapsto abcdcbcdeabcbabcdedcbcde$

6-uniform morphism $\hat{\beta}$:

а	\mapsto	bbbaa
b	\mapsto	babbaa
с	\mapsto	bbbaaa
d	\mapsto	bbaaba

 $e \mapsto bbaaab$

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Avoiding Squares in Partial Words

Binary Alphabet

Infinitely Many Holes

Theorem (B.-S., C., M. (2011))

There exists an infinite binary partial word with infinitely many holes, $w_5 = \hat{\beta}(\alpha^{\omega}(a))$, that has only the four squares: a^2 , b^2 , $(ab)^2$, $(bb)^2$.

Binary Alphabet

Complete Table

Note that $s \ge 4$, $g(h, s) = \infty$ for all non-negative integers h. Also, note that when $h \ge 5$, g(h, s) = 0 for all non-negative integers $s \le 3$.

$s \setminus h$	0	1	2	3	4	5	6
0	3	1	0	0	0	0	
1	7	5	3	0	0	0	
2	18	16	14	9	0	0	
3	∞	∞	∞	$<\infty$	$<\infty$	0	
4				∞	∞	∞	
5							

Avoiding Squares in Partial Words

Ternary Alphabet

Ternary Alphabet

What if we increase the size of the alphabet? Is it possible to get fewer squares?

Theorem (Blanchet-Sadri, Mercaș, Scott (2009))

There exist infinitely many infinite partial words with infinitely many holes over a three-letter alphabet that have at most two squares.

Is it possible to construct a partial word with one hole that has only one square?

Avoiding Squares in Partial Words

Ternary Alphabet

One Hole

The Hall morphism

$$\phi: \{a, b, c\}^* \rightarrow \{a, b, c\}^*$$

 $a \mapsto abc$
 $b \mapsto ac$
 $c \mapsto b$

Theorem

There exists an infinite square-free ternary word, $\phi^{\omega}(a)$.

Proposition (B.-S., C., M. (2011))

There exists an infinite partial word with one hole, $\diamond \phi^{\omega}(a)$, that has only one square.

Avoiding Squares in Binary Cube-free Partial Words

1 Preliminaries

Avoiding Squares in Partial Words

- Binary Alphabet
- Ternary Alphabet

3 Avoiding Squares in Binary Cube-free Partial Words

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In 1974, Entringer, Jackson, and Schatz conjecture that an infinite cube-free binary word must have an arbitrarily long square.

In 1976, Dekking shows that there exists an infinite cube-free binary word that avoids all squares xx such that $|x| \ge 4$. Dekking also proves that the bound 4 is best possible.

Note that a word that has $(aa)^2$, $(bb)^2$, $(aaa)^2$, or $(bbb)^2$ is not cubefree.

Proposition

All cubefree binary words of length greater than 113 avoiding squares xx such that $|x| \ge 4$ must have all of the following ten squares: a^2 , b^2 , $(ab)^2$, $(ba)^2$, $(aba)^2$, $(abb)^2$, $(baa)^2$, $(bab)^2$, $(bab)^2$.

Avoiding Large Squares in Partial Words — Avoiding Squares in Binary Cube-free Partial Words

Rampersad, Shallit, and Wang's Morphisms, 2

The 10-uniform morphism γ :

- $a \mapsto adbacabacd$
- $b \mapsto adbacdabac$
- $c \;\mapsto\; a c a b a d b a c d$
- $d \mapsto a cadabacab$

The 6-uniform morphism δ :

- $a \mapsto abaabb$
- $b \mapsto ababba$
- $c \mapsto abbaab$
- $d \mapsto abbaba$

Avoiding Large Squares in Partial Words — Avoiding Squares in Binary Cube-free Partial Words

Rampersad, Shallit, and Wang's Morphisms, 2

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- $a \mapsto abaabb$
- $b \mapsto ababba$
- $c \mapsto abbaab$
- $d \mapsto abbaba$

Theorem (B.-S., C., M. (2011))

There exists an infinite binary cube-free partial word with one hole, $v_1 = \diamond(\delta(a))^{-1}v_0$, that has only ten distinct squares.

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Theorem (B.-S., C., M. (2011))

Results

We found that all cube-free binary partial words of the form $u_0 \diamond u_1$ with $|u_0| = |u_1| = 9$ have at least eleven distinct squares. This implies that all cube-free binary partial words that have at most ten distinct squares must have all the holes in the first or last nine positions.

It is easy to check that it is impossible to avoid a cube in a partial word of length 11 with three holes in the first nine positions.

Results

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It is easy to check that it is impossible to avoid a cube in a partial word of length 11 with three holes in the first nine positions.

Proposition (B.-S., C., M., (2011))

All infinite cube-free binary partial words with at least three holes have at least eleven distinct squares.

Proposition (B.-S., C., M., (2011))

All cube-free binary partial words with more than four holes have at least eleven distinct squares.

Four Holes

Proposition

For large m, $v_3 = \delta(\gamma^{3m+1}(a)(acd)^{-1}) \diamond baab \diamond$, an arbitrarily long finite cube-free binary partial word with two holes, has only ten squares.

Proposition

For large m, $v_4 = \diamond baab \diamond \delta(a^{-1}\gamma^{3m+1}(a)(acd)^{-1}) \diamond baab \diamond$, an arbitraily long finite cube-free binary partial word with four holes, has only ten squares.

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Results

What happens when we have more than four holes?

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What happens when we have more than four holes?

In 1976, in the context of repetition threshold, Dejean gives an iterative morphism over a three letter alphabet that preserves $\frac{7}{4}^+$ freeness.

In 2006, Ochem gives a non-iterative morphism τ that, for any $\frac{7}{4}^+$ -free word $w \in \{a, b, c\}^*$, $\tau(w) \in \{a, b\}^*$ is $\frac{5}{2}^+$ -free, $\frac{7}{3}^+$ -free for words of length 3 and $\frac{823}{412}^+$ -free for words of length 4.

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Modification of Ochem's Morphism

We replace an occurrence of an *a* with \diamond in $\tau(b)$ where τ is the original morphism of Ochem, and get the 103-uniform morphism $\hat{\tau}$:

a → aabaabbabaababbaabaabbaabbabbabbabbaabbaabbaabbaabbaabbaabbaabbaabbabbaabbaabbabbaabbaabbabb

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Infinitely Many Holes

Theorem (B.-S., C., M. (2011))

There exists an infinite cube-free binary partial word with infinitely many holes that has only eleven distinct squares of the form xx, where $|x| \le 4$.

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Thank you!

F. Blanchet-Sadri, I. Choi, R. Mercaș, "Avoiding Large Squares in Partial Words", *Theoretical Computer Science*, 412 (2011) 3752–3758

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