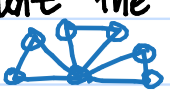


**NOTE** We need to count the zero set. Otherwise,  can be a counterexample.

**SOLUTION**

Let  $S_i$  be the number of independence set of size  $i$ . Define  $I(G; x) \equiv \sum_{i=0}^{d(G)} S_i x^i$ , ( $I(\emptyset; x) = 1$ )

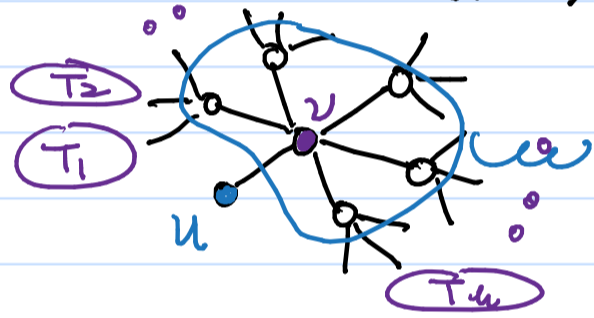
- Prop 1** For  $u, v \in V(G)$ ,  $uv \in E(G)$ ,
- 1)  $I(G; x) = I(G \setminus v; x) + x I(G \setminus v \setminus N(v); x)$
  - 2)  $I(G; x) = I(G \setminus uv; x) - x^2 I(G \setminus N(u) \setminus N(v); x)$
- $\therefore$  1) Consider independent sets with/without  $v$ .  
 2) Consider independent sets containing both  $u, v$  in  $G \setminus uv$ . I will omit the details.

**Prop 2** For a tree  $T$ ,  $I(T; -1) \in \{-1, 0, 1\}$ .  
 $\therefore$  Induction on  $n = |V(T)|$ . For  $n=1$ ,  $I(T; -1) = 0$ .  
 Suppose that Prop 2 holds for any  $k < n$  with some  $n \geq 2$ . Consider any tree with  $n$  vertices. Choose  $u \in V(T)$  with  $\deg(u) = 1$ , and let  $N(u) = \{v\}$ .

By Prop 1,

$$I(T; x) = I(T \setminus uv; x) - x^2 I(T \setminus N(u) \setminus N(v); x)$$

$$= (1+x) I(T \setminus u; x) - x^2 I(T \setminus N(v) \setminus v; x)$$



Note that  $T \setminus N(v) \setminus v = T_1 \cup \dots \cup T_k$  for some smaller trees (possibly an empty graph). Hence,  
 $I(T \setminus N(v) \setminus v; x) = \prod_{i=1}^k I(T_i; x)$ .

It follows that, from the induction hypothesis,  
 $I(T; -1) = - \prod_{i=1}^k I(T_i; -1) \in \{-1, 0, 1\}$ .

**CLAIM** Let  $G$  be a simple graph, and  $k_G$  be the smallest number of vertices required to delete every cycle. Then,  $|I(G; -1)| \leq 2^{k_G}$ .

$\therefore$  Induction on  $k_G$ . If  $k_G = 0$ , then  $G$  is a forest; thus,  $I(G; -1) \in \{-1, 0, 1\}$  by Prop 2. Suppose that  $|I(G; -1)| \leq 2^{k_G}$  for any  $G$  with  $k_G < k$  for some  $k > 0$ .

Assume that  $G$  is given and  $k_G = k$ . Let  $\{v_1, \dots, v_k\} \subseteq V(G)$  be the  $k$  vertices deleting all cycles. By the minimality, all  $v_i$  are contained in a cycle. Hence,

$$I(G; x) = I(G \setminus v_1; x) + x I(G \setminus v_1 \setminus N(v_1); x)$$

and  $k_{G \setminus v_1} = k-1$ ,  $k_{G \setminus v_1 \setminus N(v_1)} \leq k-1$ . Therefore, by the induction hypothesis,

$$|I(G; -1)| = |I(G \setminus v_1; -1) - I(G \setminus v_1 \setminus N(v_1); -1)|$$

$$\leq |I(G \setminus v_1; -1)| + |I(G \setminus v_1 \setminus N(v_1); -1)|$$

$$\leq 2^{k-1} + 2^{k-1} = 2^k,$$

and the claim follows. □