Supervised Learning Algorithms

- Gradient Descent (GD) Method

. weight update rule:

 $\underline{w}_{k+1} = \underline{w}_k - \mu \nabla_k$

where \underline{w}_{k+1} represents the k+1th n-dimensional weight vector.

true gradient (batch mode): $\nabla_k = 2R\underline{w}_k - 2P$ stochastic gradient (on-line mode): $\hat{\nabla}_k = -2\epsilon_k\underline{x}_k$

. convergence:

 $0 < \mu < 1/\lambda_{
m max}$ or $\mu \! \propto \! 1/k$

In general, this method is too slow.

. computational complexity: $O(n^2)/step$ in the batch mode O(n)/step in the on-line mode

- Newton Method

. weight update rule:

 $\underline{w}_{k+1} = \underline{w}_k - \mu R^{-1} \nabla_k$

Since $\nabla_k = 2R\underline{w}_k - 2P$ in the batch mode,

$$\begin{split} \underline{w}_{k+1} &= (1-2\mu)\underline{w}_k + 2\mu R^{-1}P = (1-2\mu)\underline{w}_k + 2\mu \underline{w}^* \text{, that is,} \\ \underline{w}_k &= \underline{w}^* + (1-2\mu)^k (\underline{w}_0 - \underline{w}^*). \end{split}$$

. convergence:

If $\mu = 1/2$, $\underline{w}_1 = \underline{w}^*$, that is, one-step convergence.

. computational complexity:

The inversion of R is required, that is,

 $O(n^3)/step$

- Quasi-Newton Method

. weight update rule in Newton method:

 $\underline{w}_{k+1} = \underline{w}_k - \alpha_k R^{-1} \nabla_k$

 $\alpha_k\,{=}\,1/2$ for one-step convergence.

. Quasi-Newton method performs the recursive construction of R^{-1} . Here, we assume that R is symmetric and positive definite matrix.

. Let us consider the following equation:

 $R\underline{p}_i = \underline{q}_i, \quad i = 0, 1, \cdots, n-1.$

Consider a matrix S_{k+1} such that

 $S_{k+1}q_i = p_i$, $i = 0, 1, \dots, k$ (Quasi-Newton condition)

After n linearly independent steps, $S_n = R^{-1}$.

construction of S_{k+1} :

$$S_{k+1} = S_k + \alpha_k \underline{r}_k \underline{r}_k^T \quad \text{(rank one correlation)} \quad \dots \quad \text{(1)}$$

$$p_k = S_{k+1} q_k = S_k q_k + \alpha_k \underline{r}_k \underline{r}_k^T q_k \quad \dots \quad \text{(2)}$$

$$q_k^T p_k = \underline{q}_k^T S_{k+1} q_k = \underline{q}_k^T S_k q_k + \alpha_k q_k^T \underline{r}_k \underline{r}_k^T q_k \quad \dots \quad \text{(3)}$$

from (2), $\alpha_k \underline{r}_k \underline{r}_k^T \underline{q}_k = \underline{p}_k - S_k \underline{q}_k$ and from (3), $\alpha_k \underline{q}_k^T \underline{r}_k \underline{r}_k^T \underline{q}_k = \underline{q}_k^T \underline{p}_k - \underline{q}_k^T S_k \underline{q}_k$. Here, $\alpha_k \underline{r}_k \underline{r}_k^T$ can be written as

$$\begin{split} \alpha_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}}^{T} &= \frac{(\alpha_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}}^{T} \underline{\boldsymbol{q}}_{\boldsymbol{k}}) (\alpha_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}}^{T} \underline{\boldsymbol{q}}_{\boldsymbol{k}})^{T}}{\alpha_{\boldsymbol{k}} \underline{\boldsymbol{q}}_{\boldsymbol{k}}^{T} \underline{\boldsymbol{r}}_{\boldsymbol{k}} \underline{\boldsymbol{r}}_{\boldsymbol{k}}^{T} \underline{\boldsymbol{q}}_{\boldsymbol{k}}} \\ &= \frac{(\underline{\boldsymbol{p}}_{\boldsymbol{k}} - S_{\boldsymbol{k}} \underline{\boldsymbol{q}}_{\boldsymbol{k}}) (\underline{\boldsymbol{p}}_{\boldsymbol{k}} - S_{\boldsymbol{k}} \underline{\boldsymbol{q}}_{\boldsymbol{k}})^{T}}{\underline{\boldsymbol{q}}_{\boldsymbol{k}}^{T} (\underline{\boldsymbol{p}}_{\boldsymbol{k}} - S_{\boldsymbol{k}} \underline{\boldsymbol{q}}_{\boldsymbol{k}})} \end{split}$$

Therefore, from (1), S_{k+1} can be written as

$$S_{k+1} = S_k + \frac{(p_k - S_k q_k)(p_k - S_k q_k)^T}{q_k^T (p_k - S_k q_k)}.$$

Let

$$\underline{y}_{k} = \frac{\underline{p}_{k} - S_{k} q_{k}}{q_{k}^{T} (\underline{p}_{k} - S_{k} q_{k})}.$$

Then,

$$S_{k+1}q_i = S_kq_i + y_k(p_k^Tq_i - q_k^TS_kq_i)$$
 for $i < k$.

By induction,

$$S_{k+1}\boldsymbol{q}_i = \boldsymbol{p}_i + \boldsymbol{y}_k (\boldsymbol{p}_k^T \boldsymbol{q}_i - \boldsymbol{q}_k^T \boldsymbol{p}_i).$$

Since $q_k^T p_i = p_k^T R p_i = p_k^T q_i$, $S_{k+1}q_i = p_i$.

This implies that after n steps,

 $S_{\!n} \boldsymbol{g}_i = \underline{p}_i \quad \text{for } 0 \leq i \leq n-1.$

Therefore, $S_n = R^{-1}$.

. Davidon-Fletcher-Powell method

Solution of rank two correlation procedure, that is,

$$S_{k+1} = S_k + \alpha_k \underline{u}_k \underline{u}_k^T + \beta_k \underline{v}_k \underline{v}_k^T.$$

The final form of S_{k+1} is described by

$$S_{k+1} = S_k + \frac{p_k p_k^T}{p_k^T q_k} - \frac{S_k q_k q_k^T S_k}{q_k^T S_k q_k}.$$

This method is more numerically stable than the rank one correlation method.

Quasi-Newton Algorithm

Step 1. Set k=0 and initialize \underline{w}_k and S_k .

(S_0 : any symmetric positive definite matrix)

Step 2. Set \underline{d}_k : $\underline{d}_k = -S_k \nabla_k$.

Step 3. Find α_k such that $\min_{\alpha_k \ge 0} E[\underline{w}_k + \alpha_k \underline{d}_k]$.

Step 4. Update $\underline{w}_{k+1}, \underline{p}_k, \underline{q}_k$:

 $\underline{w}_{k+1} = \underline{w}_k + \alpha_k \underline{d}_k, \ \underline{p}_k = \alpha_k \underline{d}_k, \text{ and } \underline{q}_k = \nabla_{k+1} - \nabla_k.$

Step 5. Update S_{k+1} : $S_{k+1} = S_k + \frac{p_k p_k^T}{p_k^T q_k} - \frac{S_k q_k q_k^T S_k}{q_k^T S_k q_k}$

Step 6. If $E[\underline{w}_{k+1}] < \theta$, stop. Otherwise, $k \leftarrow k+1$ and go to Step 2.

. convergence:

After n linearly independent steps, weight parameter can converge to \underline{w}^* .

. computational complexity: Update of S_{k+1} is required, that is,

 $O(n^2)/step$

- Lebenberg-Maquardat (LM) Method

- . The Levenberg-Marquardt algorithm was designed to approach second-order training speed without having to compute Hessian matrix.
- . The mean square error:

$$E(\underline{w}) = \frac{1}{N} \sum_{i=1}^{N} (d_i - y_i(\underline{w}))^2$$

. The Taylor series expansion of $E(\underline{w})$ around \underline{w}_{k} :

$$E(\underline{w}) \approx E(\underline{w}_k) + \nabla_k^T (\underline{w} - \underline{w}_k) + \frac{1}{2} (\underline{w} - \underline{w}_k)^T H(\underline{w}_k) (\underline{w} - \underline{w}_k)$$

where *H* is the Hessian matrix defined by $H(\underline{w}_k) \equiv [\frac{\partial E}{\partial w_i \partial w_j}]|_{\underline{w}=\underline{w}_k}$.

. Here, the Hessian matrix can be approximated as

$$H(\underline{w}_k) \approx J_k^T J_k$$

where J is the Jacobian matrix defined by

$$J_k \equiv \frac{\partial E}{\partial \underline{w}} \big|_{\underline{w} = \underline{w}_k}.$$

. The Lebenberg-Marquardt algorithm uses this approximation to the Hessian matrix:

 $\underline{w}_{k+1} = \underline{w}_k - [J_k^T J_k + \mu I]^{-1} \nabla_k.$

. μ is decreased after each successful step (reduction in performance function) and is increased only when a tentative step would increase the performance function.

. convergence:

The behavior of LM algorithm is similar to the Newton method when $\mu = 0$.

. computational complexity:

The inversion of $J_k^T J_k + \mu I$ is required, that is,

 $O(n^3)/step$

- Recursive Least Square (RLS) Method

. Let the data can be represented by $y_k = \underline{x}_k^T \underline{w}^* + n_k$

where \underline{x}_k represents the kth input vector, \underline{w}^* represents the optimal parameter vector, and n_k represents the white Gaussian noise.

. The update rule for the parameter vector:

 $\underline{w}_{k+1} = \underline{w}_k + \underline{a}_k (y_k - \underline{x}_k^T \underline{w}_k) \quad \dots \quad \text{(1)}$ where \underline{a}_k represents the gain vector. . What is the optimal gain?

determine \underline{a}_k in the sense of minimizing the mean-square distance from \underline{w}_{k+1} to \underline{w}^* .

Let

$$B_k \equiv E[(\underline{w}_{k+1} - \underline{w}^*)(\underline{w}_{k+1} - \underline{w}^*)^T]. \quad \dots \quad (2)$$

Then,

$$E[\|\underline{w}_{k+1} - \underline{w}^*\|^2] = tr \left\{ E[(\underline{w}_{k+1} - \underline{w}^*)(\underline{w}_{k+1} - \underline{w}^*)^T] \right\}, \text{ that is,}$$

 \underline{a}_k should be chosen to minimize $tr\{B_k\}$. By substituting w of (2) for w of (1) we

By substituting \underline{w}_{k+1} of (2) for \underline{w}_{k+1} of (1), we get

a new matrix B_k which depends on \underline{a}_k .

Here, the trace of B_k is determined by

$$tr\{B_k\} = tr\{B_{k-1}\} - \frac{\underline{x}_k^T B_{k-1}^2 \underline{x}_k}{1 + \underline{x}_k^T B_{k-1} \underline{x}_k} + (1 + \underline{x}_k^T B_{k-1} \underline{x}_k) \left\| \underline{a}_k - \frac{B_{k-1} \underline{x}_k}{1 + \underline{x}_k^T B_{k-1} \underline{x}_k} \right\|^2$$

The above equation is minimized when

$$\underline{a}_{k} = \frac{B_{k-1}\underline{x}_{k}}{1 + \underline{x}_{k}^{T}B_{k-1}\underline{x}_{k}}.$$

In this case,

$$B_k = B_{k-1} - \frac{(B_{k-1}\underline{x}_k)(B_{k-1}\underline{x}_k)^T}{1 + \underline{x}_k^T B_{k-1}\underline{x}_k},$$

Recursive Least Square (RLS) Algorithm

Step 1. Set \underline{w}_0 randomly, $B_0 = \epsilon I$ where ϵ is a small constant, and k = 1.

Step 2. Update
$$\underline{a}_k$$
: $\underline{a}_k = \frac{B_{k-1}\underline{x}_k}{1 + \underline{x}_k^T B_{k-1}\underline{x}_k}$

Step 3. Update B_k : $B_k = B_{k-1} - \frac{(B_{k-1}\underline{x}_k)(B_{k-1}\underline{x}_k)^T}{1 + \underline{x}_k^T B_{k-1}\underline{x}_k}$

Step 4. Update the parameter vector: $\underline{w}_{k+1} = \underline{w}_k + \underline{a}_k (y_k - \underline{x}_k^T \underline{w}_k)$ Step 5. If $tr\{B_k\} < \theta$ (threshold value), stop.

Otherwise, $k \leftarrow k+1$ and go to step 2.

. convergence:

The behavior of RLS algorithm is similar to the Quasi-Newton method. However, RLS algorithm finds \underline{w}^* asymptotically since it uses stochastic gradient $\widehat{\nabla}_k$.

. computational complexity: The calculation of B_k is required, that is,

 $O(n^2)/step$

. quick and dirty recursive linear regression: \underline{a}_{ι} is updated by

$$\underline{a}_{k} = \frac{\underline{x}_{k}}{\sum_{j=1}^{k} \|\underline{x}_{k}\|^{2}}.$$

computational complexity: O(n)/step \underline{a}_k satisfies the convergence conditions of stochastic approximation.

- Conjugate Gradient (CG) Method

Let us consider the problem of solving simultaneous equations such as

 $Q\underline{x} = \underline{b}$

This problem can be solved by minimizing the following quadratric function:

$$\min_{\underline{x}} \frac{1}{2} \underline{x}^T Q \underline{x} - \underline{b}^T \underline{x}$$

Given a symmetric matrix Q, two vectors \underline{d}_1 and \underline{d}_2 are said to be Q-orthogonal, or conjugate with respect to Qif $\underline{d}_1^T Q \underline{d}_2 = 0$ ($\underline{d}_1 \neq \underline{d}_2$). Let

$$\underline{x}^* = Q^{-1}b$$

where Q is the positive definite matrix and $\underline{d}_0, \underline{d}_1, \cdots, \underline{d}_{n-1}$ be n non-zero Q-orthogonal vectors. Then, \underline{x}^{*} can be expanded as $e^* = \alpha_0 d + \alpha_1 d_1 + \dots + \alpha_{n-1} \underline{d}$

$$\underline{x} = \alpha_0 \underline{d}_0 + \alpha_1 \underline{d}_1 + \dots + \alpha_{n-1} \underline{d}_{n-1}$$

since

$$\underline{d}_{i}^{T}Q\underline{x}^{*} = \alpha_{i}\underline{d}_{i}^{T}Q\underline{d}_{i}$$
, that is,

$$\alpha_i = \frac{\underline{d}_i^T \underline{b}}{\underline{d}_i^T Q \underline{d}_i}.$$

This implies that

$$\underline{x}^* = \sum_{i=0}^{n-1} \frac{\underline{d}_i^T \underline{b}}{\underline{d}_i^T Q \underline{d}_i} \underline{d}_i.$$

Conjugate Direction Theorem

Let $\{\underline{d}_i\}_{i=0}^{n-1}$ be a set of non-zero Q-orthogonal vectors. Then, for any $\underline{x}_0 \in \mathbb{R}^n$, the sequence $\{\underline{x}_k\}$ generated according to $\underline{x}_{k+1} = \underline{x}_k + \alpha_k \underline{d}_k, \ k \ge 0 \dots$ (1)

with

$$\alpha_k = -\frac{g_k^T \underline{d}_k}{\underline{d}_k^T Q \underline{d}_k} \text{ and } g_k = Q \underline{x}_k - \underline{b}$$

converges to the unique solution \underline{x}^* of $Q\underline{x} = \underline{b}$ after n steps, that is, $\underline{x}_n = \underline{x}^*$.

(proof)

For some set of $\alpha_k {\bf s},$

$$\begin{split} \underline{x}^* - \underline{x}_0 &= \alpha_0 \underline{d}_0 + \alpha_1 \underline{d}_2 + \dots + \alpha_{n-1} \underline{d}_{n-1} \text{ and} \\ \underline{d}_k^T Q(\underline{x}^* - \underline{x}_0) &= \alpha_k \underline{d}_k^T Q \underline{d}_k. \end{split} \text{ Then, } \alpha_k \text{ can be described by} \\ \alpha_k &= \frac{\underline{d}_k^T Q(\underline{x}^* - \underline{x}_0)}{\underline{d}_k^T Q \underline{d}_k}. \end{split}$$

By iterating (1) from \underline{x}_0 upto \underline{x}_k , we get

$$\underline{x}_k - \underline{x}_0 = \alpha_0 \underline{d}_0 + \alpha_1 \underline{d}_1 + \dots + \alpha_{k-1} \underline{d}_{k-1}.$$

From the Q-orthogonality,

$$\underline{d}_k^T Q(\underline{x}_k - \underline{x}_0) = 0, \text{ that is, } \underline{d}_k^T Q \underline{x}_k = \underline{d}_k^T Q \underline{x}_0.$$

Therefore, α_k can be redescribed by

$$\alpha_k = \frac{\underline{d}_k^T Q(\underline{x}^* - \underline{x}_0)}{\underline{d}_k^T Q \underline{d}_k} = \frac{\underline{d}_k^T Q(\underline{x}^* - \underline{x}_k)}{\underline{d}_k^T Q \underline{d}_k} = -\frac{\underline{g}_k^T \underline{d}_k}{\underline{d}_k^T Q \underline{d}_k}$$

Reference: Luenberger, "Introduction to Linear and Nonlinear Programming," chapter 8.

Conjugate Gradient Algorithm

Step 1. Set any $\underline{x}_{0} \in R^{n}$, $\underline{d}_{0} = -\underline{g}_{0} = \underline{b} - Q\underline{x}_{0}$, and k = 0. Step 2. update α_{k} : $\alpha_{k} = -\frac{\underline{g}_{k}^{T}\underline{d}_{k}}{\underline{d}_{k}^{T}Q\underline{d}_{k}}$ Step 3. update \underline{x}_{k+1} : $\underline{x}_{k+1} = \underline{x}_{k} + \alpha_{k}\underline{d}_{k}$ Step 4. update \underline{g}_{k+1} : $\underline{g}_{k+1} = Q\underline{x}_{k+1} - b$ Step 5. update β_{k} : $\beta_{k} = \frac{\underline{g}_{k+1}^{T}Q\underline{d}_{k}}{\underline{d}_{k}^{T}Q\underline{d}_{k}}$ Step 6. update \underline{d}_{k+1} : $\underline{d}_{k+1} = -\underline{g}_{k+1} + \beta_{k}\underline{d}_{k}$ Step 7. If $\|\underline{g}_{k+1}\| < \epsilon$, stop. Otherwise, $k \leftarrow k+1$ and go to step 2. . convergence:

The CG algorithm finds the unique solution \underline{x}^* of $Q\underline{x} = \underline{b}$ after n steps, that is, $\underline{x}_n = \underline{x}^*$.

. computational complexity:

The calculation of the quadratic form of $\underline{d}_k^T Q \underline{d}_k$ is required, that is, $O(n^2)/step$

. In the learning algorithm, \underline{x}_k and \underline{d}_k are associated with the weight parameter and the gradient term respectively.

- Summary of Supervised Learning Algorithms

Algorithm	Gradient Descent (GD)	Conjugate Gradient (CG)	Recursive Least Square (RLS)	Lebenberg - Marquardat (LM)	Quasi- Newton	Newton
Convergence	∞	O(n)	O(n)	1	O(n)	1
Computationa I Complexity Per Step	O(n)	O(n)	$O(n^2)$	$O(n^3)$	$O(n^2)$	$O(n^3)$
Memory Requirement	O(n)	O(n)	$O(n^2)$	$O(n^2)$	$O(n^2)$	$O(n^2)$
Learning Mode	on-line	on-line	on-line	batch	batch	batch

n: the number of parameters